

# **A Road System Interpretation of Betweenness**

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**1. Betweenness via Road Systems.** We take the intuitive view that a point  $c$  lies between points  $a$  and  $b$  exactly when every “road” allowing travel from one point to the other must go through  $c$ . I.e.,  $c$  constitutes a “complete roadblock”.

An antecedent concept that unifies most of the known examples of betweenness is that of a *road system*. This is a pair  $\langle X, \mathcal{R} \rangle$ , where  $X$  is a nonempty set and  $\mathcal{R}$  is a family of subsets of  $X$ , called *roads*, satisfying:

- Every singleton subset of  $X$  is a road. (To get from a point to itself, stay put.)
- Every doubleton subset of  $X$  is contained in at least one road. (You can always get there from here.)

Given a road system  $\langle X, \mathcal{R} \rangle$  and points  $a, b \in X$ , let  $\mathcal{R}(a, b) := \{R \in \mathcal{R} : a, b \in R\}$ ; i.e., the set of roads “joining”  $a$  and  $b$ . Then the system naturally induces a ternary (betweenness) relation  $[ , , ]_{\mathcal{R}} \subseteq X^3$ , given by stipulating  $[a, c, b]_{\mathcal{R}}$  just in case  $c \in R$  for each  $R \in \mathcal{R}(a, b)$ . The  $\mathcal{R}$ -interval  $[a, b]_{\mathcal{R}}$  is the set of points *between*  $a$  and  $b$ ; i.e.,  $[a, b]_{\mathcal{R}} := \bigcap \mathcal{R}(a, b)$ . Points  $a$  and  $b$  *bracket* the interval. (Note: an interval may have many sets of bracketing points.)

A ternary relation that is induced by a road system is called an *R-relation*.

A natural question is whether one may characterize—using first-order terms involving an abstract ternary relation symbol—exactly when a ternary relation  $[ , , ] \subseteq X^3$  is an R-relation.

This has an easy affirmative answer.

**1.1. Theorem** (Road Representation). *Let  $[ , , ]$  be a ternary relation on a nonempty set  $X$ . Then  $[ , , ]$  is an R-relation on  $X$  iff the following five first-order conditions hold:*

R1 (Symmetry)  $[a, c, b] \rightarrow [b, c, a]$   
 $([a, b] = [b, a])$ .

R2 (Reflexivity)  $[a, b, b]$   
 $([a, b] \supseteq \{a, b\})$ .

R3 (Minimality)  $[a, c, a] \rightarrow c = a$   
 $([a, a] = \{a\})$ .

R4 (Transitivity)  $([a, c, b] \wedge [a, x, c]) \rightarrow [a, x, b]$   
 $(c \in [a, b] \implies [a, c] \subseteq [a, b])$ .

R5 (Convexity)  $([a, c, b] \wedge [a, d, b] \wedge [c, x, d]) \rightarrow [a, x, b]$   
 $(c, d \in [a, b] \implies [c, d] \subseteq [a, b])$ .

R4 is an easy consequence of R5 (plus R1 and R2): just replace  $d$  with  $a$ . Why we include R4 is to point out that any ternary relation  $[ , , ]$  satisfying R1,...,R4 gives rise to a family  $\{\leq_a : a \in X\}$  of binary relations:  $x \leq_a y$  holds exactly when  $x \in [a, y]$ . Each  $\leq_a$  is reflexive (R2) and transitive (R4), and is hence a pre-order. Also the conjunction of R1, R2, and R3 says that  $a$  is the unique  $\leq_a$ -minimal element of  $X$ . We'll see later how these relations  $\leq_a$  turn out to be partial orders.

R-relations clearly satisfy R1,...,R5; and if  $[ , , ] \subseteq X^3$  satisfies R1,...,R5, we simply let  $\mathcal{R}$  be the set of intervals  $[a, b] := \{x \in X : [a, x, b] \text{ holds}\}$ . This works (but R1,...,R4 is not enough).

## 2. Classical Antecedents of Betweenness.

### 2.1. Doubly Directed Partial Orders.

A partially ordered set  $\langle X, \leq \rangle$  is *doubly directed* if each pair of elements is commonly bounded both above and below. The *D-interpretation of betweenness* in a doubly directed set is given by saying  $[a, c, b]_D$  just in case  $d \leq c \leq e$  for every  $d \leq a, b$  and  $e \geq a, b$ . By setting  $\mathcal{R}$  to comprise the order intervals  $[d, e]_O = \{x \in X : d \leq x \leq e\}$  for  $d \leq e$  in  $X$ , we immediately have  $[, , ]_{\mathcal{R}} = [, , ]_D$ .

**2.2. Trees.** A partially ordered set  $\langle X, \leq \rangle$  is a *tree ordering* if each pair of elements has a common lower bound, and no two  $\leq$ -incomparable elements have a common upper bound. The *T-interpretation of betweenness* is given by saying  $[a, c, b]_T$  holds just in case there is a common lower bound  $d \leq a, b$  such that  $d \leq c \leq a$  or  $d \leq c \leq b$ . This is clearly an R-relation; it is induced by the road system  $\mathcal{R}$ , where  $\mathcal{R}(a, b)$  comprises the sets  $V(a, b, d) := [d, a]_O \cup [d, b]_O$ ,  $d \leq a, b$ .

**2.3. Real Vector Spaces.** If  $X$  is a vector space over the reals and  $a, b \in X$ , the *S-interpretation of betweenness* is given by saying  $[a, c, b]_S$  holds just in case  $c = ta + (1 - t)b$  for some  $0 \leq t \leq 1$ . The segments  $[a, b]_S$  constitute a road system inducing  $[, , ]_S$ .

**2.4. Metric Spaces.** If  $\langle X, \rho \rangle$  is a metric space, then the *M-interpretation of betweenness* has  $[a, c, b]_M$  holding just when  $\rho(a, b) = \rho(a, c) + \rho(c, b)$ . The collection  $\mathcal{R}$ , consisting of the metric intervals  $[a, b]_M$ , is clearly a road system, but it may fail to induce  $[, , ]_M$ : K. Menger (1928) showed that, while R1, ..., R4 hold for the M-interpretation, R5 does not; it is possible for  $\bigcap \mathcal{R}(a, b)$  to be properly contained in  $[a, b]_M$ .

In a 1943 paper, M. F. Smiley considers betweenness relations that arise from multiple antecedent structures on the same underlying set. For instance, if a real vector space  $X$  has a norm  $\| \cdot \|$ , then  $X$  has an S-interpretation as well as an M-interpretation of betweenness (where  $\rho(a, b) = \|a - b\|$ ).

Betweenness in the S-sense generally implies betweenness in the M-sense, but the two notions agree exactly when the normed space is *rotund*:  $\|a + b\| = \|a\| + \|b\|$  holds for nonzero  $a, b$  iff one of the two is a non-negative scalar multiple of the other.

**3. Additive Road Systems.** A road system is *additive* if the union of two intersecting roads is a road. Additive road systems satisfy:

R6 (Disjunctivity)  $[a, x, b] \rightarrow ([a, x, c] \vee [c, x, b])$   
 $([a, b] \subseteq [a, c] \cup [c, b]).$

(To see this, suppose  $\langle X, \mathcal{R} \rangle$  is additive and  $x \notin [a, c]_{\mathcal{R}} \cup [c, b]_{\mathcal{R}}$ . Then there are  $R \in \mathcal{R}(a, c)$  and  $S \in \mathcal{R}(c, b)$  such that  $x \notin R \cup S$ . But  $c \in R \cap S$ , so  $R \cup S \in \mathcal{R}(a, b)$ . Hence  $x \notin [a, b]_{\mathcal{R}}$ .)

We also have a converse:

**3.1. Theorem** (Road Representation Continued). *An  $R$ -relation is disjunctive if and only if each of its inducing road systems is contained in an inducing road system that is additive.*

Here's what's going on with the proof:

Given a road system  $\langle X, \mathcal{R} \rangle$ , one may define  $\mathcal{R}^*$  to consist of all nonempty finite unions  $\bigcup_{i \in I} R_i$  of roads in  $\mathcal{R}$  such that whenever  $I = J \cup K$  with  $J \neq \emptyset \neq K$ , we have  $(\bigcup_{i \in J} R_i) \cap (\bigcup_{i \in K} R_i) \neq \emptyset$ . Since  $\mathcal{R}^* \supseteq \mathcal{R}$  and every superset of a road system is a road system, it's clear that the  $R$ -relation  $[\ , \ ]_{\mathcal{R}^*}$  is generally more restrictive than  $[\ , \ ]_{\mathcal{R}}$ . Indeed, the following says that  $\mathcal{R}^*$  is the *additive closure* of  $\mathcal{R}$ :

**3.2. Theorem.** *For any road system  $\mathcal{R}$ ,  $\mathcal{R}^*$  is an additive road system that contains  $\mathcal{R}$ . Moreover, if  $\mathcal{S}$  is any additive road system containing  $\mathcal{R}$ , then  $\mathcal{R}^* \subseteq \mathcal{S}$ . Finally,  $[\cdot, \cdot, \cdot]_{\mathcal{R}^*} = [\cdot, \cdot, \cdot]_{\mathcal{R}}$  iff  $[\cdot, \cdot, \cdot]_{\mathcal{R}}$  is disjunctive.*

An R-relation  $[\cdot, \cdot, \cdot]$  on  $X$  is *weakly disjunctive* if  $[a, b] \subseteq [a, c] \cup [c, b]$  for all  $c \in [a, b]$ . (The reverse inclusion is always true, so this means equality.) Here's how some of our classical antecedent structures stack up vis à vis (weak) disjunctivity:

- The D-interpretation for lattices is weakly disjunctive iff the lattice is a chain: For the “hard” direction, suppose  $a$  and  $b$  are incomparable. Then  $a \sqcup b \in [a, b]_D$ . However,  $a \sqcap b \in [a, b]_D \setminus ([a, a \sqcup b]_D \cup [a \sqcup b, b]_D)$ .
- The S-interpretation for real vector spaces is weakly disjunctive, but is disjunctive iff the vector space dimension is one: Weak disjunctivity is a simple exercise in parameterization juggling; as for non-disjunctivity when the dimension is  $\geq 2$ , it is easy to show that  $\mathcal{R}^*$ -intervals are trivial—i.e., consist of their bracket points only—whenever  $\mathcal{R}$  is a road system inducing  $[ , , ]_S$ .
- The T-interpretation for trees is disjunctive; this is not completely obvious.

**4. Separative Road Systems.** A road system  $\langle X, \mathcal{R} \rangle$  is *separative* if for any  $a, b, c \in X$  with  $b \neq c$ , there is some  $R \in \mathcal{R}$  such that either  $a, b \in R$  but  $c \notin R$  or  $a, c \in R$  but  $b \notin R$ . Recall from above the pre-orders  $\leq_a$  on  $R$ -relations; reflexive because of R2 and transitive because of R4. They become antisymmetric—and hence partial orders—in the presence of the following condition.

R7 (Antisymmetry)  $([a, b, c] \wedge [a, c, b]) \rightarrow b = c$   
 $([a, b] = [a, c] \implies b = c)$ .

**4.1. Proposition.** *A road system is separative iff its induced  $R$ -relation is antisymmetric.*

**4.2. Corollary.** *An  $R$ -relation is disjunctive and antisymmetric iff each of its inducing road systems is contained in an inducing road system that is additive and separative.*

**4.3. Proposition.** *If  $\langle X, [ , , ] \rangle$  is an  $R$ -relation that is weakly disjunctive and anti-symmetric, then for each  $a \in X$  the ordering  $\leq_a$  is a tree ordering with least element (root)  $a$ .*

In view of Proposition 4.3, each tree ordering  $\leq_a$  gives rise to a T-interpretation  $[ , , ]_a$  of betweenness on  $X$ , and it is natural to ask how these T-interpretations relate to one another and to  $[ , , ]_{\mathcal{R}}$ .

**4.4. Theorem.** *Let  $\langle X, [ , , ] \rangle$  be an antisymmetric, disjunctive  $R$ -relation, with root point  $a \in X$ . Then each interval  $[b, c]$  is contained in the corresponding tree interval  $[b, c]_a$ . Furthermore, if the centroid  $[a, b] \cap [a, c] \cap [b, c]$  is nonempty, then  $[b, c] = [b, c]_a$ .*

In weakly disjunctive, antisymmetric  $\mathcal{R}$ -relations, centroids are unique when they exist. The following example shows that the centroid existence hypothesis cannot be dropped from Theorem 4.4.

**4.5 Example.** *Let  $X$  be the unit circle, with topology inherited from the euclidean plane. Let  $\mathcal{R}$  consist of the connected subsets of  $X$ . Then  $\mathcal{R}$  is an additive, separative road system, and all  $\mathcal{R}$ -intervals are trivial. However, if  $a \in X$  is any root point, the tree order  $\leq_a$  is described by saying  $x \leq_a y$  precisely when either  $x = y$  or  $x = a$ . If  $b, c \in X \setminus \{a\}$ , then  $[b, c]_{\mathcal{R}} = \{b, c\}$ , while  $[b, c]_a = \{a, b, c\}$*

**5. Topological Antecedents of Betweenness.** There are at least three notions of betweenness in topology that deserve mention.

**5.1. The C- and Q-Interpretations** For a connected topological space  $X$ , define  $[a, c, b]_C$  (resp.,  $[a, c, b]_Q$ ) to hold if either  $c \in \{a, b\}$  or  $a$  and  $b$  lie in different components (resp., quasicomponents) of  $X \setminus \{c\}$ . The C-interpretation is induced by the additive, separative road system  $\mathcal{C}$  consisting of the connected subsets of  $X$ ; the Q-interpretation is also induced by an additive, separative road system  $\mathcal{Q}$  that is generally different from  $\mathcal{C}$ , but which may be taken to contain  $\mathcal{C}$ . (Theorems 3.1 and 3.2 come to the rescue, though proving that  $[ , , ]_Q$  satisfies R1,...,R7 takes a bit of work.) Generally the Q-interpretation is more restrictive than the C-interpretation; the two coincide in the presence of local connectedness. A point  $c \in X$  is a cut point of  $X$  precisely when it lies properly between two other points of  $X$ , in either the C- or the Q-interpretation.

**5.2. The General K-Interpretation** A connected topological space is a *continuum* if it is also compact; a *subcontinuum* of a space is a subset that is a continuum in its subspace topology. A space is *continuumwise connected* if each two of its points are contained in a subcontinuum.

We now start with a continuumwise connected topological space  $X$  and define  $[a, c, b]_K$  to hold if either  $c \in \{a, b\}$  or  $a$  and  $b$  lie in different continuum components of  $X \setminus \{c\}$ —i.e., no subcontinuum of  $X \setminus \{c\}$  contains both  $a$  and  $b$ . The K-interpretation is induced by the additive road system  $\mathcal{K}$  consisting of the subcontinua of  $X$ . In continuumwise connected spaces, the K-interpretation is less restrictive than the C-interpretation, and generally fails to be separative. A point  $c \in X$  is a *weak cut point*—i.e., its complement is not continuumwise connected—precisely when it lies properly between two other points of  $X$ , in the K-interpretation.

**5.3 The  $\mathcal{K}$ -Interpretation in Hausdorff Continua.** If  $X$  is a Hausdorff continuum, then  $\mathcal{K}$ -intervals are closed, and hence subcompacta. When they're also subcontinua, we call  $X$  *interval connected*.

For example, arcs are interval connected, as are dendrites in general. The  $\sin(\frac{1}{x})$ -continuum is another example. At the opposite extreme, in a circle (simple closed curve) any interval is trivial. Trivial intervals with at least two points are called *gaps*.

Recall that a Hausdorff continuum is *hereditarily unicoherent* if the intersection of any two of its overlapping subcontinua is a subcontinuum.

**5.3.1. Proposition.** *A Hausdorff continuum is interval connected iff it is hereditarily unicoherent.*

A primary focus in this study is characterizing being interval connected purely in first-order betweenness terms.

This issue is not yet settled, but some progress has been made.

Here are some plausible characterization sentences, listed in order of nondecreasing logical strength.

(Gap-free Property)

$$\forall a \forall b [a \neq b \rightarrow \exists c (c \in [a, b] \wedge c \neq a \wedge c \neq b)]$$

(Semi-strong Gap-free Property)

$$\forall a \forall b [a \neq b \rightarrow \exists c (c \in [a, b] \wedge c \neq a \wedge b \notin [a, c])]$$

(Strong Gap-free Property)

$$\forall a \forall b [a \neq b \rightarrow \exists c (c \in [a, b] \wedge a \notin [c, b] \wedge b \notin [a, c])]$$

The gap-free property clearly follows from interval connectedness; and, using a simple “boundary bumping” argument, we can show that the semi-strong gap-free property does as well. Not so the strong gap-free property.

**5.3.2. Theorem.** *A Hausdorff continuum satisfies the strong gap-free property iff each of its nondegenerate intervals is a decomposable subcontinuum.*

And when we strengthen gap-freeness in a completely different way, we get an even stronger condition on intervals. Recall that an R-relation is *antisymmetric* if it satisfies R7 above. In terms of intervals, this says that  $[a, b] = [a, c]$  implies  $b = c$ .

This first-order property is present in Hausdorff continua that are *aposyndetic*; i.e., any one of two points lies in the interior of a subcontinuum not containing the other.

**5.3.3. Theorem.** *A Hausdorff continuum is antisymmetric and satisfies the gap-free property iff each of its nondegenerate intervals is a generalized arc.*

So the strong gap-free property, as well as the conjunction of gap-freeness and anti-symmetry, are too strong for mere interval connectedness. On the other hand, simple gap-freeness—even semi-strong gap-freeness and the existence of centroids—is too weak.

By a *crooked annulus* we mean a continuum  $X$  that may be decomposed as a union  $X = K \cup M$ , where  $K$  and  $M$  are hereditarily indecomposable continua and  $K \cap M$  has precisely two components, each nondegenerate.

**5.3.4. Theorem.** *Every crooked annulus satisfies the semi-strong gap-free property, as well as the existence of centroids, while failing to be interval connected.*

**6. Setwise Betweenness.** Sometimes one roadblock is not enough; it may take many to prevent the bankrobbers' escape from San Diego across the border to Tijuana.

In the language of road systems, we define a *set*  $C$  to *lie between* points  $a, b \in X = \langle X, \mathcal{R} \rangle$  just in case  $C \cap R \neq \emptyset$  for all  $R \in \mathcal{R}(a, b)$ . (Note: it can still happen that  $C \cap [a, b] = \emptyset$ .)

In the setting of Hausdorff continua, we restrict attention to subsets  $C$  that are closed, indeed finite.

For  $1 \leq n < \omega$ , we denote by  $\mathcal{F}_n(X)$  the collection of nonempty subsets of  $X$  of cardinality  $\leq n$ , equipped with the Vietoris topology: For any set  $A \subseteq X$ , let  $\langle A \rangle_n^+ := \{C \in \mathcal{F}_n(X) : C \cap A \neq \emptyset\}$ ,  $\langle A \rangle_n^- := \{C \in \mathcal{F}_n(X) : C \subseteq A\}$ . The sets  $\langle U \rangle_n^+$  and  $\langle U \rangle_n^-$ , where  $U$  is open in  $X$ , constitute a subbase for a topology that makes  $\mathcal{F}_n(X)$  into a Hausdorff continuum whenever the same is true for  $X$ .

We then define  $[a, b]_n := \{C \in \mathcal{F}_n(X) : C \cap R \neq \emptyset \text{ for all } R \in \mathcal{K}(a, b)\}$ . A Hausdorff continuum is *n-interval connected* if  $[a, b]_n$  is connected for all  $a, b \in X$ .

The problem we are trying to solve is to characterize when a Hausdorff continuum is n-interval connected. Can this characterization be accomplished in first-order terms involving the obvious  $(n+2)$ -ary predicate? (We don't even know the answer for  $n = 1$ !)

Things we do know:

- For  $n \geq 2$  and  $A \subseteq X$ ,  $\langle A \rangle_n^+$  is always connected. (Proof by transfinite induction.)
- $[a, b]_n$  always contains  $\langle [a, b] \rangle_n^+$ , which is itself connected when  $n \geq 2$ . The connectedness of  $[a, b]$  is sufficient for equality to hold, but not necessary.

- Being 1-interval connected implies being  $n$ -interval connected for  $n \geq 1$ , and is the same as being hereditarily unicoherent.
- Being 2-interval connected does not imply being 1-interval connected: The unit circle is an example. (Indeed,  $[a, b]_2$  ( $a \neq b$ ) is connected, while properly containing  $\langle [a, b] \rangle_2^+ = \langle \{a, b\} \rangle_2^+$ .)
- For  $n \geq 2$ , being  $n$ -interval connected does not imply being  $(n + 1)$ -interval connected: If  $E_n$  is the union of  $n$  figure-eights joined at the tops and bottoms, then  $E_{n+1}$  is not  $(n + 1)$ -interval connected. However it is  $n$ -interval connected exactly when  $n \geq 2$ . (Indeed,  $[a, b]_n = \langle \{a, b\} \rangle_n^+$ .)

A WELSH THANKYOU:  
DIOLCH YN FAWR AM WRANDO!