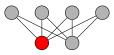


Graph theory — to the extreme!

John Engbers

Marquette University Department of Mathematics, Statistics and Computer Science

Calvin College Colloquium April 17, 2014

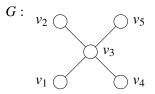


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Graph theory — to the extreme

April 2014 1 / 23

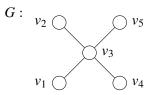
Graph theory - basics Graph G: set of vertices V(G) and a set edges E(G).



Note: Our graphs will be finite with no loops or multi-edges.

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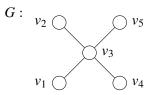
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Examples: Facebook pages/friendships; countries/border-sharing; Calvin courses/student taking both courses

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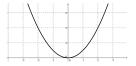


Note: Our graphs will be finite with no loops or multi-edges.

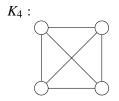
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NOT: Graph of
$$y = f(x) = x^2$$
:

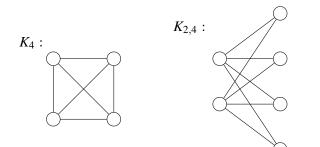


Complete graph:



Complete graph:

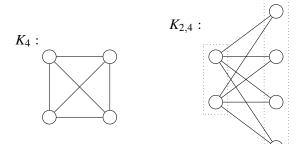
Complete bipartite graph:



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Complete graph:

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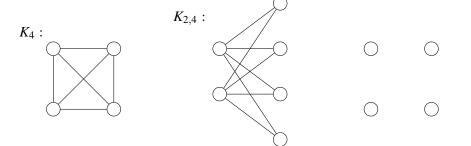
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Complete graph:

Complete bipartite graph:

Empty graph:



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Binomial coefficients: If I have n objects, then there are

$$\binom{n}{t} = \frac{n!}{t! (n-t)!}$$

different ways of selecting the *t* objects.



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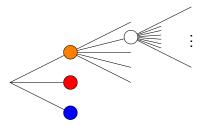
Example: Given a pool of 23 math majors, there are

$$\binom{23}{3} = \frac{23!}{3!\,20!} = 1771$$

different 3-person committees that can be formed.

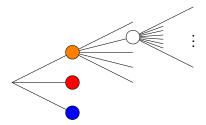


Shirts-pants-shoes idea: Suppose that I have 3 different shirts, 5 different pairs of pants, and 8 different pairs of shoes. How many different outfits can I wear?





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Answer: $3 \cdot 5 \cdot 8 = 120$ different outfits.

Extremal graph theory:

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Extremal graph theory:



Tries to figure out the "most extreme" graph from a family of graphs.

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Graph theory — to the extreme

Question (General question)

Fix a family *G* of graphs. Out of all of the graphs in *G*, which has the largest/smallest [insert something here]?

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Fix a family *G* of graphs. Out of all of the graphs in *G*, which has the largest/smallest [insert something here]?

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 $G = G_n = \{ \text{ all possible graphs on } n \text{ vertices } \}.$ [insert something here] = number of edges.

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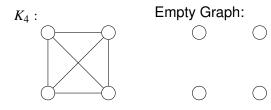
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A (10) A (10) A (10)

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Example:

 $\mathcal{G} = \mathcal{G}_{n,r} = \{ \text{ all possible graphs on } n \text{ vertices that } don't \text{ contain } K_r \}.$ [insert something here] = number of edges.

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Solution (Turán 1941): n = 7 with no K_4 :

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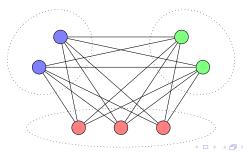
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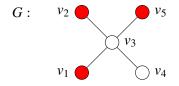
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Detour: graph theory — slightly beyond basics

Independent set (of vertices): A set of vertices which are pairwise non-adjacent.

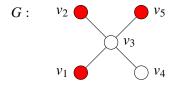


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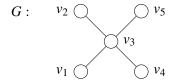


$i_t(G)$: Number of independent sets with size t in G ($t \in \{0, 1, ..., n\}$).

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In the example above,

$$i_0(G) = 1$$
, $i_1(G) = 5$, $i_2(G) = 6$, $i_3(G) = 4$, $i_4(G) = 1$, $i_5(G) = 0$

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Question (General question)

Fix a family *G* of graphs. Out of all of the graphs in *G*, which has the largest/smallest [insert something here]?

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 $\mathcal{G} = \mathcal{G}_n = \{ \text{ all graphs on } n \text{ vertices } \}.$ [insert something here] = $i_t(G)$ for each t = 0, 1, 2, ..., n

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Empty Graph: $K_4 :$

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Question (Today's question - different family)

Fix a family \mathcal{G} of graphs. Out of all the graphs in \mathcal{G} , which has the largest value of $i_t(G)$ for each t?

Intuition: fewer edges implies more independent sets of a fixed size.

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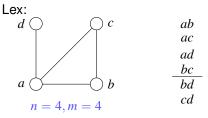
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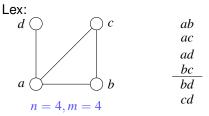


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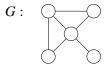
- $\mathcal{G} = \{ all graphs with$ *n*vertices and*m* $edges \}.$
 - (Cutler, Radcliffe 2011) Lex graph maximizes $i_t(G)$ for all t.



Is there a more indirect (local) way to force lots of edges?

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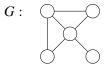
Minimum degree δ : Smallest number of edges adjacent to a vertex



 $\mathcal{G}_n(\delta) = \{ \text{All graphs on } n \text{ vertices with minimum degree at least } \delta \}.$

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Minimum degree δ : Smallest number of edges adjacent to a vertex

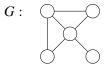


 $\mathcal{G}_n(\delta) = \{ \text{All graphs on } n \text{ vertices with minimum degree at least } \delta \}.$

Question (Today's question)

Out of all the graphs in $\mathcal{G}_n(\delta)$, which has the largest value of $i_t(G)$ for each *t*?

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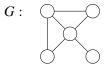
Remarks (true for *all* graphs on *n* vertices):

• For all graphs *G*, $i_0(G) = 1, i_1(G) = n$.

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Extremal graph theory

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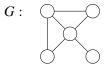
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Guess: extremal graph in $\mathcal{G}_n(\delta)$ should have all degrees equal to δ .

 $\delta = 0, 1$

$\mathcal{G}_n(0) \implies$ empty graph maximizes $i_t(G)$. \checkmark

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Question (Today's question, $\delta = 1$)

Out of all the graphs in $\mathcal{G}_n(1)$, which has the largest value of $i_t(G)$ for each *t*?

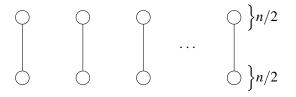
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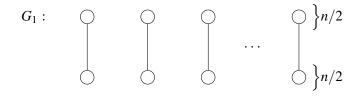
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Out of all the graphs in $\mathcal{G}_n(1)$, which has the largest value of $i_t(G)$ for each *t*?

Extremal guess (graph with the least number of edges, *n* even):

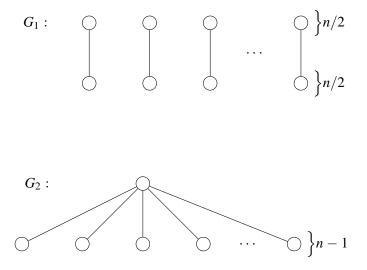


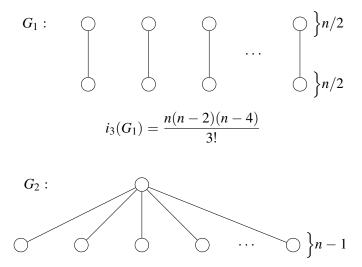
Has largest value of $i_t(G)$ for t = 0, 1, 2 and for all G in $\mathcal{G}_n(1)$.



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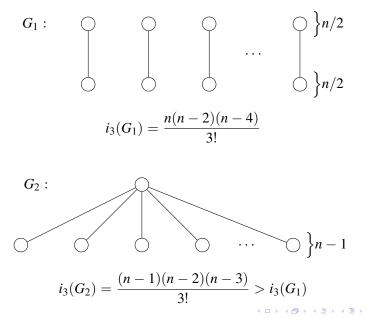
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John Engbers (Marquette University)

April 2014 14 / 23



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Old thought: extremal graph has fewest number of edges.

New thought: extremal graph has largest maximal independent set (for all $t \ge 3$).

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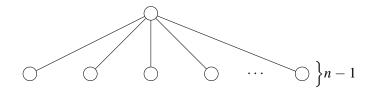
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Theorem (Galvin, 2011)

For each $3 \le t \le n-1$, any graph $G \in \mathcal{G}_n(1)$ has

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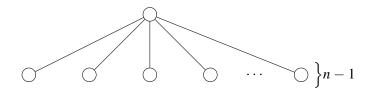
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Let's start by looking at size t = 3: Let *G* be any graph with minimal degree at least 1.

• Suppose that *G* has a vertex *v* as pictured:



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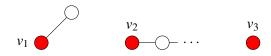


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- **Goal:** show if we don't have this situation, then we have at most $\binom{n-1}{3}$ independent sets of size 3.

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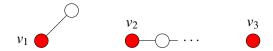
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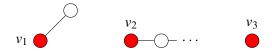
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How many ordered independent sets of size 3 can G have? At most:

$$n(n-2)(n-4) < (n-1)(n-2)(n-3).$$

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So G has at most

$$\frac{n(n-2)(n-4)}{3!} < \frac{(n-1)(n-2)(n-3)}{3!} = \binom{n-1}{3}$$

unordered independent sets of size 3.

We can use ordered independent sets to obtain the result for any t > 3:

How many ordered independent sets of size 4 can *G* in $\mathcal{G}_n(1)$ have?

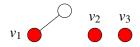
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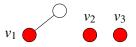
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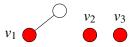
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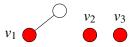


- There are at most (n-1)(n-2)(n-3)(n-4) ordered independent sets of size 4.
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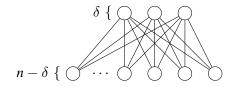
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- This argument works for any $t \ge 4$.

General Result

Larger minimum degrees? Tweaking the previous argument gives:

Theorem (E., Galvin 2014) Fix $\delta \ge 2$, size $\delta + 1 \le t \le n - \delta$, and *n* large enough. If $G \in \mathcal{G}_n(\delta)$, then

$$i_t(G) \leq i_t(K_{\delta,n-\delta}) = \binom{n-\delta}{t}$$



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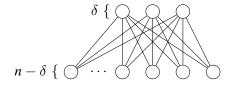
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Conjecture

Let $G \in \mathcal{G}_n(\delta)$ for $\delta > 2$. Then for each size $3 \le t \le \delta$,

$$i_t(G) \leq i_t(K_{\delta,n-\delta}) = \binom{n-\delta}{t} + \binom{\delta}{t}.$$



Hot off the presses:

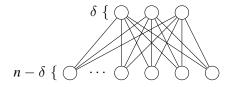
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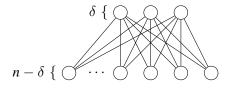
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- Cutler-Radcliffe [2014] showed $\sum_{t} i_t(G) \leq \sum_{t} i_t(K_{\delta,n-\delta})$
- Gan-Loh-Sudakov [2014+] Conjecture true(!) (for all $n \ge 2\delta$)



What if...?

John Engbers (Marquette University)

Graph theory — to the extreme

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What if ...?

"In mathematics, the art of proposing a question must be held of higher value than solving it." — Georg Cantor



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Question: What if $n < 2\delta$? Example: n = 7, $\delta = 4$

John Engbers (Marquette University)

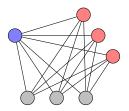
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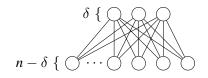
Question: What if $n < 2\delta$? Example: n = 7, $\delta = 4$



Does this graph maximize $i_t(G)$ ($t \ge 3$) when $n < 2\delta$? [True if $n - \delta | \delta |_{\Omega \otimes \Omega}$

John Engbers (Marquette University)

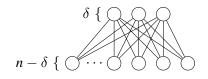
Graph theory — to the extreme



Other open questions:

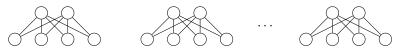
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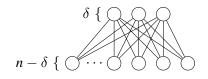
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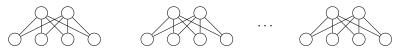
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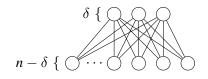


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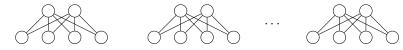


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. . .

• Question: Special case: What if all vertices have degree δ ?

• (Kahn) Conjectured extremal graph (for δ -regular graphs):

Thank you!

Slides available on my website: www.mscs.mu.edu/~engbers/

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